

## A PRELIMINARY STUDY OF A FUZZY CLUSTERING AND ASSIGNMENT PROBLEM-BASED CELL FORMATION ALGORITHM

Susanto, S., Kennedy, R.D., and Price, J.W.H.

Department of Mechanical Engineering, Monash University – Australia

(E-mail: [ssusanto@home.unpar.ac.id](mailto:ssusanto@home.unpar.ac.id))

### ABSTRACT:

The ideal situation of CM is achieved when each part-cluster together with the machine-cluster for the manufacturing cell is independent from the rest of manufacturing cell, that is exceptional elements are non-existent. However, such a situation is rarely achieved. In this paper, a cell formation algorithm that identifies part-clusters and machine-clusters separately is introduced. The algorithm is based on fuzzy clustering and the assignment technique. Computational experience shows that this new algorithm can give superior performance, especially for higher density machine-part incidence matrices.

### KEYWORDS:

Cellular Manufacturing, Manufacturing Cell, Cell Formation, Part-cluster, Machine-cluster.

## 1. INTRODUCTION

GT in particular is a manufacturing-related philosophy that is gaining [1]. This philosophy exploits the similarity among the attributes of given objects [2]. One of the popular applications of GT is called *cellular manufacturing* (CM). In CM, parts that require similar machines are grouped into the same *part-cluster*, while machines needed by a part-cluster are grouped to form a *machine-cluster*. A *manufacturing cell* is part of a manufacturing system where a machine-cluster is placed to process a part-cluster. The most distinctive feature of CM is that it contains dissimilar machines that are located in one area.

The first problem which must be addressed when considering a cellular manufacturing system is that of *cell formation*. Data for cell formation is organized into a *machine-part incidence matrix*. This is a binary matrix with 0 or 1 entries. A 1(one) entry in row  $i$  and column  $j$  of the matrix indicates that machine-type  $i$  needs to operate on part-type  $j$ , while a 0(zero) indicates it does not. The cell formation problem in CM is the partition or decomposition of the manufacturing system into (manufacturing) cells. Mathematically, the problem involves converting the machine-part incidence matrix into a block diagonal form in which 1(one) entries are concentrated in blocks along the diagonal of the matrix. Each block represents a manufacturing cell [3].

Extensive work has been done by many researchers to provide new techniques for solving this problem. The cell formation problem is complicated by the existence of exceptional parts and/or exceptional machines [4]. An *exceptional part* is a part that requires processing in another machine-cluster. An *exceptional machine* is a machine that processes parts from a different part-cluster. Both exceptional parts and exceptional machines cause *intercellular movement* of parts. Ideally a part-cluster is processed in a single machine cell for its entire operations. In practice, however, it is a very rare case [5]. Thus, the more 1(one) entries are concentrated in the block diagonal the more effective a cell formation technique is. Numerous research papers have appeared in the literature for cell formation. These methods are based on the following approaches: coding and classifications, machine-component analysis, similarity coefficients, mathematical programming and heuristic methods, knowledge-based and pattern recognition methods, neural networks and fuzzy clustering [2].

It is the last approach (ie. fuzzy clustering) that we focus on in this article. Very few papers have appeared in the area of fuzzy clustering techniques applied to cell formation. The application of fuzzy clustering techniques in cellular manufacturing originally proposed by Chu and Hayya [6] and was followed by Ponnambalam and Aravindan [7] with a minor modification in the initialization. Despite their contributions in introducing the application of this technique, there is still a room for improvement, especially in the effort of increasing the 1(one) entries to be concentrated in the block diagonal of the initial given machine-part incidence matrix. The main feature of this paper is to demonstrate a new algorithm to effect this improvement as outlined in Section 3. In Section 2, the basic techniques of the algorithm are discussed. Those techniques are *the fuzzy c-means* (fuzzy clustering) and *the assignment problem*. In Section 4 some examples of computational results are presented. A summary and suggestions for further research are presented in Section 5.

## 2. A BRIEF REVIEW

### 2.1 Fuzzy Clustering Problem

Clustering is a process which is common and basic to human understanding. This process, which plays a vital role in grouping of related objects, can be found in such diverse fields as statistics, economics, physics, psychology, biology, pattern recognition, engineering, marketing [8,9]. The primary objective of *clustering* is to partition a given set of objects into so-called homogeneous clusters (groups) [10-12]. Globally, the clustering

problem can be divided into two main categories ie. *hard clustering*, in which an object belongs only to one cluster, and *fuzzy clustering*, in which every object belongs to all clusters with different degrees of membership. Definitions, theorems and algorithms involved in the discussion of the fuzzy clustering topic are as follows [13].

### 2.1.1 Definition and Theorem

#### Definition 1. Fuzzy c-Partitions

Let  $O = \{o_1, o_2, \dots, o_p\}$  be a set of  $p$  objects;  $X = \{x_1, x_2, \dots, x_p\}$  where  $x_r = (x_{1r}, x_{2r}, \dots, x_{mr})^T$  is the attribute vector of object  $o_r \in O$  ( $r = 1, 2, \dots, p$ ); and  $V_{pc}$  is the set of real  $p \times c$  matrices;

Fuzzy c-Partitions space for  $X$  is the set

$$M_{fc} = \left\{ U \in V_{pc} \mid u_{ri} \in [0,1] \forall r,i; \sum_{i=1}^c u_{ri} = 1 \forall r; 0 < \sum_{r=1}^p u_{ri} < p \forall i \right\} \quad (2.1)$$

#### Definition 2. Fuzzy c-Means Functional

Fuzzy c-Means Functional is the functional  $J_f : M_{fc} \times R^{pc} \rightarrow R^+$  defined by the relationship

$$J_f(U, v) = \sum_{r=1}^p \sum_{i=1}^c (u_{ri})^f (d_{ri})^2 \quad (2.2)$$

where

- 1)  $U \in M_{fc}$  is a fuzzy c-partitions of  $X$ ;
- 2)  $v = (v_1, v_2, \dots, v_c) \in R^{mc}$  with  $v_i = (v_{1i}, v_{2i}, \dots, v_{mi})^T \in R^m$  is the cluster centre of  $u_i = (u_{1i}, u_{2i}, \dots, u_{pi})^T$ ,  $1 \leq i \leq c$ ;
- 3)  $d_{ri} = \|x_r - v_i\|$  and  $\|\cdot\|$  is any inner product on  $R^m$
- 4)  $f$  is the weighting exponent (also called *the degree of fuzziness*),  $f \in (1, \infty)$ .

#### Theorem 1. (Bezdek's Necessary Conditions for Global Minimum)

Let  $\|\cdot\|$  be any inner product norm on  $R^h$ ;  $f \in (1, \infty)$ ;  $X$  has at least  $c$  ( $c < p$ ) distinct points, and for each  $r \in \{1, 2, \dots, p\}$ , define the sets

$$I_r = \{i \mid 1 \leq i \leq c; d_{ri} = \|x_r - v_i\| = 0\} \quad (2.3)$$

$$I_r^c = \{1, 2, \dots, c\} - I_r \quad (2.4)$$

then if  $(U, v) \in M_{fc} \times R^{pc}$  is the global minimum point for  $J_f$  then

$$1) I_r = \emptyset \Rightarrow u_{ri} = \frac{1}{\left[ \sum_{k=1}^c \left( \frac{d_{ri}}{d_{rk}} \right)^{\frac{2}{f-1}} \right]}, \quad (2.5)$$

$$\text{or} \\ I_r = \emptyset \Rightarrow u_{ri} = 0 \forall i \in I_r^c \text{ and } \sum_{i \in I_r} u_{ri} = 1; \quad (2.6)$$

$$2) v_i = \frac{\sum_{r=1}^p (u_{ri})^f x_r}{\sum_{r=1}^p (u_{ri})^f} \quad \forall i \in \{1, 2, \dots, c\} \quad (2.7)$$

### 2.1.2 Chu and Hayya's (1991) Algorithm

Chu and Hayya [6] were the first authors to apply the fuzzy clustering technique to cellformation. They have treated the attribute vector of all part type as an input to their algorithm. Their algorithm will result in two different matrices. The first matrix, notated as  $U$ , *the degree of membership matrix*, is used for grouping parts into part-clusters. The second matrix, notated as  $v$ , *the cluster centre matrix*, is used for grouping machines into machine-clusters. The detail of their algorithm is as follows.

Let  $X = (x_{qr})_{m \times p}$  be the given machine-part incidence matrix of  $m$  types of machine and  $p$  types of parts;  $x_r = (x_{1r}, x_{2r}, \dots, x_{mr})^T$ , namely the  $r$ -th column of  $X$  ( $r = 1, 2, \dots, p$ ) is the *attribute vector* of the part type  $r$ . The  $p$  part types and the  $m$  machine types are to be grouped into  $c$  clusters.

**Algorithm 1. Chu and Hayya's Algorithm for part-cluster and machine-cluster formation**

**Step 1:** Fix  $c$ ,  $2 \leq c < \min\{m, p\}$ ; Choose any positive value  $\xi$  for stopping criterion;. Choose any inner product norm  $\| \cdot \|$  on  $R^m$  and fix  $f \in (1, \infty)$ ; initialize  $U^{(0)} \in M_c$ . Set  $l = 0$ .

**Step 2:** Calculate the  $c$  fuzzy cluster centres  $\{v_{(i)}^{(l)}\}$ , where

$$v_{qi}^{(l)} = \frac{\sum_{r=1}^p (u_{ri}^{(l)})^f x_{qr}}{\sum_{r=1}^p (u_{ri}^{(l)})^f}; \quad i=1,2,\dots,c; \quad q=1,2,\dots,m \quad (2.8)$$

**Step 3:**  $l \leftarrow l + 1$ ; calculate

$$u_{ri}^{(l)} = \begin{cases} 1 & \text{if } I_r = \emptyset \\ \left[ \frac{1}{\sum_{k=1}^c \left( \frac{d_{ri}}{d_{rk}} \right)^{\frac{2}{f-1}}} \right] & \text{if } I_r \neq \emptyset \\ 0, \forall k \in I_r^c & \text{if } I_r \neq \emptyset \\ \frac{1}{|I_r|} & \text{if } I_r \neq \emptyset \end{cases} \quad (2.9)$$

where  $I_r$  and  $I_r^c$  are as defined in (2.4) and (2.5) respectively,  $|I_r|$  is the number of element(s) in  $I_r$ .

**Step 4:** If  $\|U^{(l+1)} - U^{(l)}\| < \xi$  then stop else goto Step 2.

The final matrices  $U$  and  $v$  are used to determine the part-clusters and machine-clusters respectively, under the following rule:

**Rule 1. (Part-clusters formation Procedure)**

If  $u_{ri} = \max_k \{u_{rk}\}$ ;  $r = 1, 2, \dots, p$  ( $p$  is the number of part-types) and  $k = 1, 2, \dots, c$  ( $c$  is the number of part clusters), then part  $r$  is assigned to part-cluster  $i$ .

**Rule 2. (Machine-clusters formation Procedure)**

If  $v_{qi} = \max_k \{v_{qk}\}$ ;  $q = 1, 2, \dots, m$  ( $m$  is the number of machine-types) and  $k = 1, 2, \dots, c$  ( $c$  is the number of machine-clusters) then machine  $q$  is assigned to machine-cluster  $i$ .

As a result of Chu and Hayya's stopping criterion, two problems may arise. The *first problem*, there is a distinct possibility that their algorithm results in some empty part-clusters and/or empty machine-clusters (ex. see result displayed in Table 3. in [7]). Furthermore, their algorithm may end up with different numbers of part-clusters compared to the resulting number of machine-clusters, which in turn, does not have any practical meaning. The *second problem*, in case of a tie in Rule-1, that is  $\max_k \{u_{rk}\}$  is achieved by more than one part-cluster, then in Chu and Hayya's algorithm [6] (or its minor modification by Ponnambalam and Aravindan [7]) the part-type  $r$  is directly assigned to the first part-cluster achieving the maximum value of  $u_{rk}$ . This is too restrictive in the authors' opinion, since it eliminates the opportunity for the other part-cluster that achieved the same maximum value of  $u_{rk}$  to contain that particular part type. Improvements to overcome these problems is discussed in Section 3.

## 2.2 The Assignment Problem

Consider the situation of assigning  $c$  part-clusters to  $c$  machine-clusters and define  $c_{ik} = \sum_{r \in \text{PC}-k} \sum_{q \in \text{MC}-i} x_{qr}$  be the degree of conformance of part-cluster(PC)  $k$  and machine-cluster (MC)  $i$ . That is, the total number of '1'(one) entries in the block diagonal formed by those part and machine-clusters. The objective is to assign the  $c$  part-clusters to the  $c$  machine-clusters (one part-cluster per machine-cluster) to maximize the total degree of conformance. This situation is known as the assignment problem. (see also [14=15]).

The assignment problem, if applied in cell formation situations, can be expressed in mathematical model as follows. Let

$$\Omega_{ij} = \begin{cases} 0, & \text{if the part-cluster(PC) } i \text{ is not assigned to the machine-cluster(MC) } j \\ 1, & \text{if the part-cluster(PC) } i \text{ is assigned to the machine-cluster(MC) } j \end{cases} \quad (2.10)$$

The model is thus given by

$$\text{maximize } z = \sum_{i=1}^c \sum_{j=1}^c c_{ij} \Omega_{ij} \quad (2.11)$$

subject to

$$\sum_{j=1}^c \Omega_{ij} = 1, \quad i = 1, 2, \dots, c \quad (2.12)$$

$$\sum_{i=1}^c \Omega_{ij} = 1, \quad j = 1, 2, \dots, c \quad (2.13)$$

$$\Omega_{ij} = 0 \text{ or } 1. \quad (2.14)$$

Having obtained the part-clusters and the machine-clusters, the application of the assignment problem to the cell formation problem will maximize the total number of elements in the block diagonals of the final permuted machine-part incidence matrix. Thus, at the same time, this will minimize the incidence of exceptional parts and machines.

## 3. THE SKP-1 (Susanto-Kennedy-Price Version 1) ALGORITHM

The SKP-1 algorithm is aimed at tackling the two problems identified in Chu and Hayya's algorithm [6] and at maximizing the total number of elements in the block diagonal of the final permuted machine-part incidence matrix. As classified by Wemmerlöv [16] (in [17]), the cell formation literature can be divided into four categories, according to the formation logic used: 1) grouping part-clusters only; 2) forming part-clusters and then machine-clusters or vice versa; 3) forming part-clusters and machine-clusters simultaneously, and 4) grouping machine-clusters only.

The SKP-1 algorithm belongs to the third category in the forementioned classification, and can be outlined as follows. Let  $X = (x_{qr})_{m \times p}$  be the given machine-part incidence matrix of  $m$  types of machine and  $p$  types of parts;  $x_r = (x_{1r}, x_{2r}, \dots, x_{mr})^T$ , namely the  $r$ -th column of  $X$  ( $r = 1, 2, \dots, p$ ) is the attribute vector of the part type  $r$ . Let  $y_q = (x_{q1}, x_{q2}, \dots, x_{qp})$ , namely the  $q$ -th row of  $X$  ( $q = 1, 2, \dots, m$ ) is the attribute vector of the machine type  $q$ . The  $p$  part types and the  $m$  machine types are to be grouped into  $c$  clusters ( $2 \leq c < \min\{m, p\}$ ) separately using the fuzzy clustering approach.

### 3.1 The First Step of SKP-1 (Machine-clusters Formation)

The  $p$  attribute vectors of each part type (ie.  $x_1, x_2, \dots, x_p$ ) serve as the only known variable in the functional to be minimized  $J_f(U, v)$  (as defined in (2.2)). In this minimization problem, we define a successful machine-cluster formation solution as one that 1) converges to  $c$  non empty machine-clusters; 2) results in unimodal solution for  $\max_k \{v_{qk}\}$  ( $q = 1, 2, \dots, m$  and  $k = 1, 2, \dots, c$ ).

Instead of using the criterion in Step-4 of Algorithm-1, we will use a given number of maximum iterations. Thus, up to the maximum number of iterations, it is possible to get more than one successful machine-cluster formation solution. The details of the first step in SKP-1 are as follows:

**Step 1** : Fix  $c$ ,  $2 \leq c < \min\{m, p\}$ ; choose any inner product norm  $\| \cdot \|$  on  $R^m$  and fix  $f \in (1, \infty)$ ; initialize  $U^{(0)} \in M_{fc}$ . let  $success = 0$ ;  $l = 0$  and the number of maximum iterations be  $max\_iterations$ .

**Step 2** : Calculate the  $c$  fuzzy cluster centres  $\{v_{(i)}^{(l)}\}$  using (2.8).

**Step 3** : If  $\forall q \in \{1, 2, \dots, m\}$ , the set  $A_q = \left\{ \alpha_l^{(q)} \mid v_{q\alpha_l^{(k)}} = \max_{i=1,2,\dots,c} \{v_{qi}^{(l)}\} \right\}$  is singleton

then if  $\bigcup_{q=1}^m A_q = \{1, 2, \dots, c\}$

then

if  $l < max\_iterations$

then 1)  $success \leftarrow success + 1$  2) at the  $l$ -th iteration and at the  $success$ -th successful machine assignment, assign machine  $q$  to machine-cluster  $\alpha_l^{(q)}$

else 1)  $success \leftarrow success + 1$  2) at the  $l$ -th iteration and at the  $success$ -th successful machine assignment, assign machine  $q$  to machine-cluster  $\alpha_l^{(q)}$  3) goto Step 6

else if  $l = max\_iterations$  then goto Step 6

else if  $l = max\_iterations$  then goto Step 6

**Step 4** : Determine  $I_r$  and  $I_r^c$  as defined in (2.3) and (2.4), where  $r = 1, 2, \dots, p$

**Step 5** :  $l \leftarrow l + 1$ ; calculate  $\{u_{fi}^{(l)}\}$  as defined in (2.9), goto Step 2.

**Step 6** : If  $success \geq 1$  then  $MC = \bigcup_{s=1}^{success} MC_s$ , where  $MC_s = \{MC_{is} \mid i = 1, 2, \dots, c; s = 1, 2, \dots, success\}$  and

$MC_{is}$  is the machine-cluster  $i$  obtained from the  $s$ -th successful machine assignment, else "this algorithm fails to converge to a successful machine-cluster formation solution".

### 3.2 Second Step of SKP-1 (Part-clusters Formation)

Basically the second step of SKP-1 is congruent to that of the first step. While the first step of SKP-1 exploits the part-attribute vectors  $x_r$  ( $r = 1, 2, \dots, p$ ) (which are column-vectors) to form the machine-clusters, the second step exploits the machine-attribute vector  $y_q$  ( $q = 1, 2, \dots, m$ ) (which are row-vectors) to form the part-clusters. We define a *successful part-cluster formation solution* as one that 1) converges to  $c$  non-empty part-clusters and 2) results in a unimodal solution for  $\max_k \{v_{qk}\}$  ( $k = 1, 2, \dots, c$ ).

Again, instead of using the stopping criterion in Step-4 of Algorithm-1 (Section 2.1.2) we will use a given number of maximum iterations. Thus, up to the maximum number of iterations, it is possible to get more than *one successful part-cluster formation solution*. If, after the maximum number of iterations achieved, this second step converges and results in *SUCCESS successful part-cluster formation solutions* ( $SUCCESS \geq 1$ ), then let us define:

$PC = \bigcup_{S=1}^{SUCCESS} PC_S$ , where  $PC_S = \{PC_{kS} \mid k = 1, 2, \dots, c; S = 1, 2, \dots, SUCCESS\}$  and  $PC_{kS}$  is the part-cluster

$k$  obtained from the  $S$ -th successful part-cluster assignment.

### 3.3 The Third Step of SKP-1 (The Assignment Problem)

To apply this step, it is necessary that both the first and the second step of SKP-1 converge to *success* successful machine-cluster solutions and *SUCCESS* successful part-cluster solutions. Let

$$C_{MC_{is}PC_{kS}} = \sum_{r \in PC_{kS}} \sum_{q \in MC_{is}} x_{qr} \quad (2.15)$$

be the *degree of conformance* of part cluster  $k$  obtained from the  $S$ -th successful part-cluster formation and machine cluster  $i$  obtained from the  $s$ -th successful machine-cluster formation.

The final assignment problem is the solution of the following optimization problem:

$$\text{maximize } z_{SS} = \max \sum_{i=1}^c \sum_{k=1}^c C_{MC_{is}PC_{kS}} \Omega_{MC_{is}PC_{kS}} \quad (2.16)$$

subject to

$$\sum_{k=1}^c \Omega_{MC_i, PC_{kS}} = 1, i=1,2,\dots,c; s=1,2,\dots, success; S=1,2,\dots, SUCCESS \quad (2.17)$$

$$\sum_{i=1}^c \Omega_{MC_i, PC_{kS}} = 1, \quad k = 1,2,\dots,c; s=1,2,\dots, success; S=1,2,\dots, SUCCESS \quad (2.18)$$

$$\Omega_{MC_i, PC_{kS}} = 0 \text{ or } 1 \quad (2.19)$$

**Note:**  $\Omega_{MC_i, PC_{kS}} = 1$ , if the  $k$ -th part-cluster (obtained from the  $S$ -th part-cluster formation solution) is assigned to the  $i$ -th machine-cluster (obtained from the  $s$ -th machine-cluster formation solution), and  $= 0$ , otherwise.

#### 4. COMPUTATIONAL RESULTS

In this section the performances of SKP-1 and Chu and Hayya's algorithm are compared, based on the observations applied to 16 matrices with the dimension of 40x50, with density ranges from 20 to 80%. The results are presented in Table 1.

**Table-1. Comparison of percentage of non exceptional elements resulted from SKP-1 and Chu and Hayya's(1991) algorithm**

MATRIX	Density of Matrix *)	SKP -1	CHU & HAYYA [6]	Comment on SKP-1
<b>20 – 40%</b>				
I	0.207	37.7	35.7	better
II	0.236	38.3	31.6	better
III	0.254	32.3	34.5	worse
IV	0.260	36.0	42.5	worse
V	0.295	25.8	37.4	worse
VI	0.352	29.6	30.7	worse
<b>40 – 60%</b>				
VII	0.404	31.4	30.4	better
VIII	0.431	28.3	26.0	better
IX	0.445	28.2	29.4	worse
X	0.459	31.8	24.8	better
XI	0.485	35.5	25.6	better
<b>60 – 80%</b>				
XII	0.611	26.1	25.6	better
XIII	0.659	33.4	29.5	better
XIV	0.669	29.5	23.0	better
XIII	0.704	28.8	24.8	better
XV	0.730	24.4	30.2	worse
XVI	0.805	25.6	20.7	better

\*) Density of a matrix as defined in [19]

#### 4.1 Example of Computational Result

The machine-part incidence matrix XVI is displayed in Appendix 1. The results of using this matrix to Chu and Hayya's [6] and the SKP-1 algorithm are outlined in the following Section.

##### 4.1.1 Result from Chu and Hayya's algorithm [6]

This algorithm was applied to matrix XVI with  $c = 6, f = 2, \xi = 0.001$  and Euclidean norm as the inner product. The machine-clusters and part-clusters are as follows.

**The Machine-clusters :**

MC - 1 = {4,7,15,17,35,37}; MC - 2 = {6,19,20,21,23,24,30,31,38,40}; MC - 3 = {8,10,36}; MC - 4 = {3,22}; MC - 5 = {1,5,9,11,12,14,25,26,27,32,33,39} and MC - 6 = {2,13,16,18,28,29,34}.

**The Part-clusters :**

PC-1 = {1,7,18,19,20,25,33,37,43,49}; PC-2 = {14,32,38,50}; PC-3 = {5,9,21,31,39,40,45}; PC-4 = {4,10, 22, 28,44,46,47}; PC-5 = {2,11,15,17,23,24,26,27,29,34,35,41} and PC-6 = {3,6,8,12,13,16,30,36,42,48}

**The resulted Manufacturing Cells are :**

Manufacturing Cell-*i* will consist of MC-*i* that processes PC-*i* (*i* = 1,2, ..., 6).

There will be 333 of '1'(one) entries on the block diagonal of the final rearrangement of the matrix XVI which contains 1609 (ie.  $\approx 0.805 \times 40 \times 50$ ) '1' entries. Thus, Chu and Hayya's [6] algorithm results in 20.7% (ie.  $333/1609 \times 100\%$ ) of non exceptional elements. The final rearrangement by this algorithm is displayed in the Appendix 2.

**4.1.2 Result from the SKP-1 algorithm**

The SKP-1 algorithm was also applied to matrix XVI,  $c = 6$ ,  $f = 2$ ,  $\xi = 0.001$ ,  $max\_iterations = 200$ , and Euclidean norm as the inner product. Four successful machine-cluster formation solutions and five successful part-cluster formation solutions resulted from the treatment are:

**The Four Successful Machine-cluster Formation Solutions :**

*The first successful machine-cluster formation solution:*

MC - 1 = {1,2,3,4,7,15,17,35,37}; MC - 2 = {6,19,23,25,30,31,38,40}; MC - 3 = {8,10,18,21,24,36}; MC - 4 = {20,22,27}; MC - 5 = {5,9,11,14,26,32,33,39} and MC - 6 = {12,13,16,28,29,34}.

*The second successful machine-cluster formation solution:*

MC - 1 = {1,3,4,7,15,16,17,35,36,37,40}; MC - 2 = {6,19,23,31,38}; MC - 3 = {8,10,21,24,30}; ;  
MC - 4 = {20,22}; MC - 5 = {5,9,11,14,25,26,27,32,33,39} and MC - 6 = {2,12,13,18,28,29,34}

*The third successful machine-cluster formation solution:*

MC - 1 = {15,16,17,20,35,36,37,38,40}; MC - 2 = {19,31}; MC - 3 = {7,8,10,12,24,30}; MC - 4 = {3};  
MC - 5 = {1,5,9,14,21,22,25,26,27,32,33,39} and MC - 6 = {2,4,6,11,13,18,23,28,29,34}

*The fourth successful machine-cluster formation solution:*

MC-1 = {15,16,18,20,31,35,36,37,38,40}; MC-2 = {19}; MC - 3 = {7,8,10,12,17,24,30}; MC-4 = {3}  
MC - 5 = {1,5,9,14,21,22,25,26,27,32,33,39} and MC - 6 = {2,4,6,11,13,23,28,29,34}

**The Five Successful Part-cluster Formation Solutions :**

*The first successful part-cluster formation solution:*

PC- 1 = {6,7,13,22,25,32}; PC- 2 = {34,35,46,47,49}; PC- 3 = {10,14,15,17,26,27,30,31,43,44,45}; PC- 4 = {1,8,12,16,18,19,20,21,23,28,33,36,37,38,48,50}; PC- 5 = {11} and PC- 6 = {2,3,4,5,9,24,29, 39,40,41,42}

*The second successful part-cluster formation solution:*

PC- 1 = {6,13,18,25,32}; PC- 2 = {34,35,47,49}; PC- 3 = {10,14,15,17,26,27,31,43,45}; PC- 4 = {1,7,8,12,16, 19, 20,21,22,23,28,33,36,37,38,42,48,50}; PC- 5 = {11,30,46} and PC- 6 = {2,3,4,5,9,24,29,39,40,41,44}

*The third successful part-cluster formation solution:*

PC- 1 = {4,6,7,13,18,22,25,32,38}; PC- 2 = {34,49}; PC- 3 = {10,14,15,17,26,27,29,31,43,45,47}; PC- 4 = {1,8,12,16,19,20,21,23,28,33,35,36,37,42,48,50}; PC- 5 = {11,30,41,46}; PC- 6 = {2,3,5,9,24,39,40,44}

*The fourth successful part-cluster formation solution:*

PC- 1 = {4,7,13,18,22,32,38}; PC- 2 = {3,34,49}; PC- 3 = {6,10,14,15,17,25,26,27,29,31,43,47}; PC- 4 = {1,8, 12,16,19,20,21,23,28,33,35,36,37,39,42,48,50}; PC- 5 = {2,11,24,30,41,45,46} and PC- 6 = {5,9,40,44}

*The fifth successful part-cluster formation solution:*

PC- 1 = {4,7,18,22,32,38}; PC- 2 = {49}; PC- 3 = {6,9,10,14,15,17,25,26,27,29,31,34,43,47}; PC- 4 = {1,8,12, 16,19,20,21,23,28,33,35,36,37,39,42,48,50}; PC- 5 = {2,3,11,13,24,30,41,45,46} and PC- 6 = {5,40,44}

The desired machine-clusters and part-clusters are obtained by pairing each successful machine-cluster and part-cluster formation solution, and applying the Assignment Algorithm to each pair. In the case of machine-part incidence matrix XVI, the desired machine-clusters are those obtained from the *fourth* successful machine-cluster formation solution and the desired part-clusters are those obtained from the *fourth* successful part-cluster formation problem.

**The resulted Manufacturing Cells are :**

M-Cell 1 = (MC-1, PC-4); M-Cell 2 = (MC-2, PC-6); M-Cell 3 = (MC-3, PC-1); M-Cell 4 = (MC-4, PC-2); M-Cell 5 = (MC-5, PC-3) and M-Cell 6 = (MC-6, PC-5).

There will be 412 of 1(one) entries on the block diagonal of the final rearrangement of the matrix XVI out of 1609 entries. Thus, this algorithm results in 25.6% of non exceptional elements. The final rearrangement by this algorithm is displayed in the Appendix 3..

## 5. SUMMARY AND SUGGESTIONS

In general, computational experience summarized in Table 1 shows that for higher density machine-part incidence matrix, the proposed new algorithm, ie. Susanto-Kennedy-Price Version 1 or SKP-1, has more chance in giving cell formation solutions with fewer numbers of exceptional elements.

However, further research is required to identify other attributes of the input machine-part incidence matrix, such as *the total bond energy* matrix (as defined in [19]), in causing exceptional elements in the final solution.

## ACKNOWLEDGMENT

We would like to thank the Australian Agency for International Development (AusAID), formerly AIDAB (Australian International Development Assistance Bureau) for their support of this research, and the Monash University Research Training and Support Branch for providing additional funding. Also, special acknowledgment is directed toward Dr. P.A. Grossman, Department of Mathematics, Monash University, for his time in the informal discussions and personal correspondence relating to Fuzzy Clustering and the Assignment Problem.

## REFERENCES

1. Wemmerlöv, U., Production Planning and Control Procedures for Cellular Manufacturing Systems, The American Production and Inventory Control Society, Inc., Falls Church, USA, 1988.
2. Singh, N., "Design of cellular manufacturing systems: An invited review", *European Journal Of Operational Research*, Vol. 69, p. 284-291, 1993.
3. Seifoddini, H., and Djassemi, M., "The threshold value of a quality index for formation of cellular manufacturing systems", *International Journal of Production Research*, Vol. 34, No. 12, p. 3401-3416, 1996.
4. Boe, J.W., and Cheng, C.H., "A close neighbour algorithm for designing cellular manufacturing systems", *International Journal of Production Research*, Vol. 29, No. 10, p. 2097-2116, 1991.
5. King, J.R., "Machine-component grouping in production flow analysis: An approach using rank order clustering algorithm", *International Journal of Production Research*, Vol. 18, p. 213-232, 1980.
6. Chu, C.H., and Hayya, J.C., "A fuzzy clustering approach to manufacturing cell formation" *International Journal of Production Research*, Vol. 29, p. 1475-1497, 1991.
7. Ponnambalam, S.G., and Aravindan, P., "Design of cellular manufacturing systems using objective functional clustering algorithms", *International Journal of Advanced Manufacturing Technology*, Vol. 9, p. 390-397, 1994.
8. Brown, D.E., and Huntley, C.L., "A practical application of simulated annealing to clustering", *Pattern Recognition*, Vol. 25, No. 4, p. 401-412, 1992.
9. Zhang, Q, and Boyle, R.D., "A new clustering algorithm with multiple runs of iterative procedures", *Pattern Recognition*, Vol. 24, No. 9, p. 835-848, 1991.
10. Ruspini, E., "A new approach to clustering", *Information Control*, Vol. 15, No. 1, p. 22-32, 1969.
11. Gitman, I., and Levine, M., "An algorithm for detecting unimodal fuzzy sets and its application as a clustering technique", *IEEE Trans. Comput.*, Vol. C-19, No. 7, p. 583-593, 1970.
12. Bezdek, J.C., and Pal, S.K., Fuzzy Models for Pattern Recognition, IEEE Press, Piscataway, NJ, USA, 1992.
13. Bezdek, J.C., Pattern Recognition with Fuzzy Objective Function Algorithms, Plenum Press, New York, 1981.
14. Taha, H., Operations Research: an introduction, 5th. ed., p. 214, Macmillan Publishing Company, Don Mills, Ontario, 1992.
15. Bunday, B.D., Basic Linear Programming, Edward Arnold, London; Baltimore; Md, USA, 1984.
16. Wemmerlöv, U., and Hyer, N.L., "Research issues in cellular manufacturing", *International Journal of Production Research*, Vol. 25, 413-431, 1987.
17. Chu, C.H., "Recent advances in mathematical programming for cell formation", Planning, Design, and Analysis of Cellular Manufacturing Systems, ed. Kamrani, A., Parsaei, H.R., and Liles, D.H., p. 3-46, Elsevier Science B.B., Amsterdam, 1995.
18. Miltenburg, J., and Zhang, W., "A comparative evaluation of nine well-known algorithms for solving the cell formation problem in group technology", Journal of Operations Management, Vol. 10, No. 1, January, 1991.
19. McCormick, W.T, Schweitzer, R.J, and White, T.W., "Problem decomposition and data reorganization by clustering techniques", *Operations Research*, Vol. 20, p. 993-1009, 1972.

APPENDICES

Appendix 1. The Machine-Part Incidence Matrix XVI

		Part-type Number																																							
		000000000111111111122222222233333333334444444445										12345678901234567890123456789012345678901234567890										1234567890123456789012345678901234567890																			
01	Machine Type Number	1111111111011111011110111101111010111101111111111																																							
02		111101110111111111111111101111011111111111111001110111111																																							
03		0001111011011110111101001111111110111111111111111011																																							
04		1100101110111101111110111110111011000110101111110010111																																							
05		0101110101101110110001110111111111111010101111111111																																							
06		11110011111110100111111010011111010011110100111110111111																																							
07		11101101111011111111111111010111101011011011101010																																							
08		111010111101111111011011011110111111011111111111100010																																							
09		111100111100111110000111111111000111111100111111001100011																																							
10		101111111010111010111111110011100011111011111111111110																																							
11		11111010101111011011101111101111111011111111111011111111																																							
12		01111101100110111111101101101110011101111111110001110																																							
13		111111100111101111111011110111101111111111111110110101																																							
14		0111111011111111011111111111111111111111111001001011011110																																							
15		101111101111110111111111101111111101101101101111101111																																							
16		101011111001110111111110011101111011111110111111010010111																																							
17		101110111111011011111111111111111111111111111110011001111																																							
18		01111101110111110111110111111111111111111111110111111101101																																							
19		100110111111111101111110111111111111011111111111111110101																																							
20		111111111111111001111110011111111111110101010101011111																																							
21		111111111111111110111011110111111111111111111111111111111																																							
22		1100110111111111111011101111111111111111111111111101111101																																							
23		11101001111111010111111110010101110111111111111111111111																																							
24		11111111111111011111011011110111100011111111111111110111																																							
25		0100111110101111111110111111111111111111111111011101111011																																							
26		111011010111111110101111110101101111111111111111101111111																																							
27		01111011111111111101111101111011001111011111111111111101																																							
28		11111111001111010111111111111111111111111111111111111110111																																							
29		111101110111110110110011101111100111111111111111111111101																																							
30		111110011111011100111011111111011011111111111111111111111																																							
31		011110111101110001111111111110110111111111111111111111111																																							
32		11101011111111111110001111111111111111110011101101110111110																																							
33		111001000111111111111111111110111101100111111111111111111																																							
34		1011011111111001011111110101111111111111111111111111111111																																							
35		101111111111111111111111111011111111111101101111111110101110																																							
36		1111111111011111111111101110101111111111111111111111111011																																							
37		11010111110111111111111111111011111111111111111111111111111																																							
38		10110111011101100111001101101111111111111111111111111011111																																							
39		101101111110111010001010011111111011110101111111111111110																																							
40		1111001100110111111111111001011111101111111111111111011111																																							

Appendix 2. The Final Rearrangement of Machine-Part Incidence Matrix XVI by Chu and Hayya's Algorithm

		Part-type Number																																							
		0011223344 1335 0023344 0 122444										011122222334 0001113344										1789053739 4 280 5911905 4 028467 215734679451 3682360628																			
4	Machine Type Number	1111111111 1111 1110110 0001010 110111010011 0011110011																																							
7		1111111111 1110 0111011 1111101 11111100011 1101011000																																							
15		1111100111 1111 1111011 1111101 010111111110 1101111011																																							
17		1111111111 1111 1111110 1111101 011111111110 1011001101																																							
35		1111110111 1110 1111011 1111001 011101111111 1111110111																																							
37		1111111111 1111 0111011 1111111 101110111110 0111110111																																							
6		1001101001 1111 1111111 1110111 110011101011 1011111011																																							
19		1111101110 1111 1111111 1111110 011011111011 0011111111																																							
20		1101101001 1111 1111010 1111111 111011011110 1111111111																																							
21		1110111111 1111 1111111 1110111 111101111111 1111111011																																							
23		1011111111 1111 1110111 0111111 110011000101 1011111111																																							
24		1111101011 1111 1111111 1101111 111111111101 1111100010																																							
30		1001110111 1111 1111111 1101111 111011111111 1011010111																																							
31		0111111101 1111 0111111 1111111 100011101111 1111100111																																							
38		1111101111 1111 0001110 1100111 011011111111 1111001011																																							
40		1111100111 1111 0011110 1011111 111110101111 1011011111																																							
8		1110101011 1110 1111111 0101100 101111111111 1011110110																																							
10		1110110111 1110 1111111 1110111 001011110000 1111011111																																							
36		1111111111 1111 1111111 1100111 101111011111 1111111110																																							
3		0111111111 1111 1101111 1111111 001100111101 0101101110																																							
22		1010111110 1111 1111011 0111111 111101111111 0111111101																																							
1		1101110111 1111 1111101 1110111 111101111101 1110111111																																							
5		0010001111 1101 1001101 1111111 111111111111 0110101011																																							
9		1010010101 1011 1101111 1101100 111111111011 1010011100																																							
11		1101111111 1101 1111111 1000111 110111111111 1001111111																																							
12		0011100011 0010 1111110 1000001 101111111111 1111111111																																							
14		0111111011 1110 1111000 1111111 111111111111 1101101001																																							
25		0111110011 1111 1101111 0011111 111111111110 0110111110																																							
26		1001011101 1011 1011111 0111111 111111100111 1111111111																																							
27		0110100110 1001 1111111 1111111 111111110111 1011111111																																							
32		1100010011 1010 1111101 0111011 111111111111 1011111111																																							
33		1011110011 1111 0011111 0110111 111111111111 1101111011																																							
39		1100000011 1110 0111011 1101111 011110111111 1110101111																																							
2		1111111111 1111 0011001 1110011 111101111111 1111111111																																							
13		1111110110 1111 1011111 1010010 110101111111 1111111111																																							
16		1111101101 1111 1111111 0011001 000110110010 1111111111																																							
18		0011111110 1111 1111111 1111101 101001111111 1111111011																																							
28		1111111111 1111 1010101 1011110 110011111111 1111111111																																							
29		1101111110 1011 0000101 1101111 110111011111 1111111111																																							
34		1111101111 0111 0111101 1111111 010011101111 1111111111																																							

**Appendix 3. The Final Rearrangement of Machine-Part Incidence Matrix XVI by Susanto-Kennedy-Price Version 1 (SKP-1) Algorithm**

	Part-type Number											
	00111222233333445	044	001233	4	00111122223344	001123444	18269013835679280	504	478228	9	69045756791437	231340156
15	10111111101010111	111	111111	1	111101011111111	011111010						
16	11111111111111111	110	011111	1	110101011101001	010101010						
18	01111110111011111	111	101111	0	111110111111111	110111110						
20	11111111111100111	111	110111	1	111110001111101	111111001						
31	01101111111111111	011	111111	1	111100110111101	110110111						
35	11111110101110110	110	111111	1	111111111111111	011110110						
36	11111110111111101	111	111011	1	111111101111111	110111111						
37	11111111111110111	011	111111	1	111111111111111	100100011						
38	11101101011011111	011	111011	1	101110011111111	011011101						
40	11111111101111111	011	111111	1	000111010111111	111001101						
19	11111111111111111	111	111111	0	011110011111010	001111111						
7	10111111111010000	011	111111	1	11111111001011	111011110						
8	11110111111101100	111	011011	1	011111011111110	110110110						
10	11110111000111110	111	111111	1	111110111101011	010011011						
12	01111111001101110	110	101001	1	110011011111111	110111100						
17	11101111111111011	111	111111	1	011111111111111	011011000						
24	11101111110001101	111	111011	1	111111011111111	111110111						
30	11111111101111111	111	100011	1	011111011111111	111010111						
Machine Type Number	3	00101100110111101	111	111111	1	111111111111111	000101111					
1	11011110000111111	101	110111	1	111111111111111	111111111						
5	01000001111011111	101	101110	1	101111011111111	101111111						
9	11010001101111001	111	101001	1	01111111111000	111011110						
14	00101111111000010	101	111111	1	111111111111111	111111101						
21	11110110011011111	111	111111	1	111111111111111	111111111						
22	11110110111111001	111	001111	0	111111111111111	101111111						
25	01011101101101101	111	011111	1	110111111111111	101111011						
26	11111011111111111	111	000101	1	10111111001101	111111111						
27	01110111101111111	111	111100	0	011111011011111	111111111						
32	11110011101101110	100	010101	1	011111111111111	111111111						
33	10111110010011111	011	001111	1	101111111111111	111111111						
39	11000011101100110	011	110011	1	111111011111111	011101111						
2	11111110011110111	000	111111	1	101111111111111	111111111						
4	11111111111011111	110	011011	1	01010110100010	101110101						
6	11111111011001111	111	100111	1	01110001011001	111111111						
11	10111111011111111	111	110010	1	010101111111111	111111111						
13	11111110001111111	110	111111	0	100101111111110	111111111						
23	11111111110111111	111	001111	1	01110010000111	111111111						
28	11111111111111111	101	111111	1	10010011110110	111111111						
29	11111101111111111	001	110001	0	10110110110111	111111111						
34	11111111111111111	001	111111	1	11100001011111	011111111						